## Automation in the Maintenance of the Linux Kernel The Coccinelle Experience

Julia Lawall, Gilles Muller (Inria/LIP6) February 6, 2019

## What is an operating system kernel?



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#### Contains:

- screen
- keyboard
- camera
- CPU
- disk
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# The OS kernel is the software that manages application access to hardware

## Operating system correctness is critical

- Our "modern" society (almost) entirely relies on computers
- Applications are directly impacted by any bug in the operating system
  - Application unavailability
  - Loss of data
  - Security attacks
  - Risks for human life
  - Impossible to predict the next exploit



#### Our focus: The Linux kernel

- Open source OS kernel, developed by Linus Torvalds
- First released in 1991
- Version 1.0.0 released in 1994
- Today used in the top 500 supercomputers, billions of smartphones (Android), battleships, stock exchanges, ...



#### Development challenges

- Large code size
  - 17.5 million lines of code in Linux v4.20 (Dec 2018)
- Multiple streams of development
  - Mainline version, accepting new features and bug fixes.
  - Stable versions, accepting only bug fixes from the mainline.
- Wide range of contributors
  - Almost 20 000 contributors since 2005.
  - Industry developers, hobbyists, newbies.
- High rate of change and heavy code review burden
  - 13-14K commits (changes) per release (every 2-3 months).
  - Need to be integrated with both mainline and stable versions.

Code size



7

Release model



A new mainline release every 2-3 months. 13-14K commits (changes) per release

## Maintainer hierarchy (tiny extract)



1287 maintainers in v4.20 (Dec 2018), each responsible for 1-8000 files.

#### Collaborative development

Largest contributors:  $\geq$  10 developers and  $\geq$  500 commits per year.



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## Contribution workflow

#### All contributions go through email



#### Emailed patch, fixing a bug in a drm driver

From nobody Sun Jan 13 09:05:41 CET 2019

From: Julia Lawall <Julia.Lawall@lip6.fr>

- To: Sandy Huang <hjc@rock-chips.com>
- Cc: "Heiko Stübner" <heiko@sntech.de>,David Airlie <airlied@linux.ie>,Daniel Vetter <daniel@ffwll.ch>, dri-devel@lists.freedesktop.org,linux-arm-kernel@lists.infradead.org,linux-rockchip@lists.infradead.org, linux-kernel@vger.kernel.org

Subject: [PATCH 3/4] drm/rockchip: add missing of\_node\_put

```
The device node iterators perform an of_node_get on each iteration, so a jump out of the loop requires an of_node_put.
```

```
Signed-off-by: Julia Lawall <Julia.Lawall@lip6.fr>
---
drivers/gpu/drm/rockchip/rockchip_rgb.c | 4 +++-
1 file changed, 3 insertions(+), 1 deletion(-)
```

```
diff --git a/drivers/gpu/drm/rockchip/rockchip_rgb.c b/drivers/gpu/drm/rockchip/rockchip_rgb.c
index 96ac145..37f9302 100644
a contract of the second s
```

```
--- a/drivers/gpu/drm/rockchip/rockchip_rgb.c
```

+++ b/drivers/gpu/drm/rockchip/rockchip\_rgb.c

```
aa -116,2 +116,4 aa struct rockchip_rgb *rockchip_rgb_init(struct device *dev,
```

- if (!ret)

4

```
+ if (!ret) {
```

```
+ of_node_put(endpoint);
```

break;

## Problem: Big code implies the need for big changes

Thomas Gleixner (maintainer of x86 architecture support, interrupt support, etc.): Remove the irq argument from interrupt handlers.

• 188 affected files.

Kees Cook (organizer of the Linux kernel self-protection project): Replace malloc-based array allocations by safer versions.

• 377 and 484 affected files.

Deepa Dinamani (developer), and many others: Time data size for Y2038.

 $\cdot~\sim$  250 affected files.

Julia Lawall: Add missing of\_node\_put.

• 50 affected files.

## Some patches fixing missing of \_node\_puts

#### drivers/video/backlight/88pm860x\_bl.c

+ of\_node\_put(np);
 break;

#### drivers/power/charger-manager.c

#### arch/arm/kernel/devtree.c

```
രു -101,6 +101.7 രു
 if (of_property_read_u32(cpu, "reg", &hwid)) {
        pr debug(" * %s missing reg property\n",
                     cpu->full name):
        of node put(cpu):
        return;
aa -108.8 +109.10 aa
   * 8 MSBs must be set to 0 in the DT since the reg
   * defines the MPIDR[23:0].
   */
- if (hwid & ~MPIDR HWID BITMASK)
+ if (hwid & ~MPIDR HWID BITMASK) {
       of_node_put(cpu);
        return:
+ }
 /*
```

\* Duplicate MPIDRs are a recipe for disaster.

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        return:
+ }
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 * Duplicate MPIDRs are a recipe for disaster.
```

Challenge: How to make all these changes quickly, consistently and correctly?

- Allows changes to C code to be expressed using patch-like code patterns (semantic patches) using the language SmPL.
- Applies SmPL semantic patches to an entire code base, updating all relevant code sites at once.
- Under development since 2005. Open source since 2008.

- Like patches, but independent of irrelevant details (line numbers, spacing, variable names, etc.)
- Derived from code, with abstraction.
- Easily adaptable, to eliminate false positives or treat new cases.
- Goal: fit with the existing habits of the Linux programmer.

- Like patches, but independent of irrelevant details (line numbers, spacing, variable names, etc.)
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#### No program analysis expertise required.

## Example: Find and fix missing of\_node\_put bugs

#### Memory management in kernel code.

- **kmalloc** to allocate memory.
- kfree later, to free it.
- Requires knowing when a free is possible, i.e., no remaining pointers.

## Example: Find and fix missing of\_node\_put bugs

#### Memory management in kernel code.

- **kmalloc** to allocate memory.
- kfree later, to free it.
- Requires knowing when a free is possible, i.e., no remaining pointers.

#### Some parts of the kernel provide more support.

- Reference counting for device nodes.
- of\_node\_get on access.
- of\_node\_put to allow freeing.

#### of\_node\_get's and of\_node\_put's clutter the code.

• Bookkeeping, not functionality.

Solution? Hide them when possible.

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Each iteration increases the reference count of a new element, and decreases the reference count of the previous element.

## for\_each\_child\_of\_node hides reference counts in the normal case.

- Problem: Abnormal loop exits don't benefit from the hidden puts:
  - break, return, goto
  - Not continue
- Consequence: memory leak.

#### Evolution of the problem over time



#### There are currently 9 such iterators

• for\_each\_of\_cpu\_node added in August 2018.

#### An example to motivate semantic patch construction

```
for each child of node(port, endpoint) {
     if (of property read u32(endpoint, "reg", &endpoint id))
         endpoint id = 0;
      if (rockchip drm endpoint is subdriver(endpoint) > 0)
         continue:
      child count++;
      ret = drm_of_find_panel_or_bridge(dev->of_node, 0, endpoint_id,
                                        &panel, &bridge);
     if (!ret)
     if (!ret) {
+
         of node put(endpoint):
         break;
```

Check execution paths through and after the loop body.

```
an
expression node,child;
iterator name for_each_child_of_node;
an
for_each_child_of_node(node, child) {
   ...
}
...
```

## All ok if reach an **of\_node\_put** or add a pointer to **child**. **break** requires further attention.

```
രി
expression node, child, e;
iterator name for_each_child_of_node;
രി
for each child of node(node, child) {
  . . .
  of node put(child):
  e = child
  break;
  . . .
```

. . .

```
Add of_node_put before break;.
```

```
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expression node, child, e;
iterator name for_each_child_of_node;
രി
for each child of node(node, child) {
  . . .
  of_node_put(child);
  e = child
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  break:
  . . .
```

. . .

#### Make **break** optional.

. . .

```
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for each child of node(node, child) {
  . . .
  of_node_put(child);
  e = child
+ of_node_put(child);
? break;
  . . .
```

Allow any code to follow the matched code in the loop. By default, the code matched by a pattern adjacent to ... cannot be matched in ...

```
രി
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iterator name for_each_child_of_node;
രി
for each child of node(node, child) {
  . . .
  of node put(child):
  e = child
+ of_node_put(child);
? break:
  ... when any
```

Check that **child** is not used after the loop.

```
രി
expression node, child, e;
iterator name for_each_child_of_node;
രി
for each child of node(node, child) {
  . . .
  of_node_put(child);
  e = child
+ of_node_put(child);
? break:
  ... when any
... when != child
```

## Results: Identified missing of\_node\_puts in Linux 4.20 (Dec 2018)

	break	return	goto
for_each_child_of_node	4	71	20

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	break	return	goto
for_each_child_of_node	4	71	20
for_each_node_by_name	1	1	0
for_each_node_by_type	1	1	0
for_each_compatible_node	2	7	0
for_each_matching_node	2	3	1
for_each_matching_node_and_match	1	8	1
for_each_available_child_of_node	19	23	18
for_each_node_with_property	0	0	0
for_each_of_cpu_node	0	0	0

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	break	return	goto
for_each_child_of_node	4	71	20
for_each_node_by_name	1	1	0
for_each_node_by_type	1	1	0
<pre>for_each_compatible_node</pre>	2	7	0
for_each_matching_node	2	3	1
for_each_matching_node_and_match	1	8	1
for_each_available_child_of_node	19	23	18
for_each_node_with_property	0	0	0
for_each_of_cpu_node	0	0	0

Studying these results reveals other types of reference count problems for which we can write semantic patches.

#### Impact: Coccinelle-related Linux kernel commits per year





#### Impact: Intel's 0-day test service



- How to write semantic patches?
- How to validate the results?
- How to submit the resulting patches?
- Scaling to multiple development streams.

#### How to write semantic patches?

#### Kernel developers are C programmers, not SmPL programmers.

- Some use Coccinelle regularly, and become experts
  - Thomas Gleixner, Kees Cook, etc.
- Others make widespread changes infrequently.
  - Learn/relearn SmPL on their own.
  - Ask us to solve the problem (of\_node\_put).
  - Ask another developer to solve the problem.
  - Make the change without Coccinelle, and make a note to do better next time.

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Spinfer: Infer a semantic patch from a few change examples (in progress).

Coccinelle makes it easy to make lots of changes, very fast.

- Changes needed validation, whether or not generated using Coccinelle.
- Testing is hard, may need unavailable hardware, specific inputs.
- Even compilation is hard, due to configuration options.

#### make coccicheck

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0-Day, from Intel: compilation for many configurations and build testing.

- Runs a selection of Coccinelle scripts, via make coccicheck.
- Almost 600 0-day reports motivated by Coccinelle rules in 2017.

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JMake: Make for kernel janitors [DSN 2017].

• Chooses a relevant configuration and reports whether changed code was subjected to the compiler.

Coccinelle makes it easy to make lots of changes, very fast.

· Changes need to go to the right person in the maintainer hierarchy.



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Coccinelle has been suggested as a solution to synchronize treewide changes with the current kernel version.

#### Challenge: Scaling to multiple development streams



- Backports project uses Coccinelle to retarget version specific code to a version generic library [EDCC 2015]
- Prequel finds commits that illustrate how to port drivers across versions [USENIX ATC 2017]
- Spinfer infers semantic patches from such examples

#### Conclusion

- Coccinelle, provides user scriptable matching and transformation of C code.
- Lessons learned:
  - Take the expertise of the target users into account.
  - Avoid creeping featurism: Do one thing and do it well.
- Success measures:
  - Over 6600 commits in the Linux kernel based on Coccinelle.
  - EuroSys test of time award (2018).

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## http://coccinelle.lip6.fr