## Quantum Computing as a Service

## Secure and Verifiable Multi=Tenant Quantum Data Centre



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## Quantum Computing and Simulation Hub

CNRS Sorbonne University Quantum Internet Alliance

VeriQloud


## Currently

## Quantum Links



## Quantum Links

Unclonable / Measurement disturbance ... - security QKD, Quantum Coin Flipping, ...

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Unclonable / Measurement disturbance ... - security QKD, Quantum Coin Flipping, ...

Shangıal


## Quantum Nodes



## Quantum Links

Unclonable / Measurement disturbance ... - security QKD, Quantum Coin Flipping, ...

## Quantum Nodes

Superposition / Entanglement... - speed


Random Walk, Machine Learning, ...

## Quantum Links



Future

## Multi-Tenant Quantum Data Centre



## Multi-Tenant Quantum Data Centre



## Use-Case Example: Privacy Preserving QML



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Party with Data


## Quantum Secure Multi Party Computing

Party with Q Computer

Plan

Plan

- 2 party QC: Honest Client - Malicious Server


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- What is possible?
- Building Blocks: QKD, Teleportation, Self-Testing
- Verifiable Universal Blind Quantum Computing


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- Multi party QC: Malicious Clients - Malicious Server
- Lifting Classical SMPC to QSMPC
- When can we have it for real?


## Honest Client - Malicious Server



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## Honest Client - Malicious Server



## Secure Cloud Computing

Rivest, Adleman and Dertouzos 1979
Can we process encrypted data without decrypting it first ?

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## Limited Client

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## Untrusted Server

## Secure Cloud Computing

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Rivest, Adleman and Dertouzos 1979
Can we process encrypted data without decrypting it first ?


Gentry 2009 - Fully Homomorphic Encryption
computational security computational security

## Secure Classical access to Quantum Cloud?

Fillinger: No efficient informationally secure classical FHE scheme exist

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Murimae: No informationally secure quantum scheme for quantum function evaluation (for restricted classical client)

## Secure Classical access to Quantum Cloud?

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On the implausibility of informationally secure quantum cloud computing with Classical Client (PH collapses at the third level)

Aaronson, Cojocaru, Gheorghiu, Kashefi, 2017

## Generalised Encryption Scheme (GES)



Which functions admit a GES?


## What about NP functions?



Unless PH collapses

## Generalised Encryption Scheme for QC (GES)



## Our work

1. Do BQP functions admit a GES?

We give evidence that the answer is NO


Conjectured relationship between classes

## An oracle result

For each d, there exists an oracle, O, such that:


The oracle is based on Simon's problem

$$
O(n, x)=f_{n}(x)
$$

Is $f_{n} 1$-to- 1 or does it have Simon's property?
Simon's property: $f_{n}$ is 2-to- 1 and periodic

## A sampling result



Unless, there exist circuits $\left\{C_{n}\right\}_{n}$ having the properties:

$$
\begin{gathered}
\left|C_{n}\right|=2^{n-\Omega(n / \log (n))} \\
C_{n} \text { queries } \mathrm{NP}^{\mathrm{NP}}
\end{gathered}
$$

Computes exactly the permanent of $n \times n$ matrix
Best known algorithm for permanent (Ryser '63): $O\left(n 2^{n}\right)$

## A sampling result



GES for SampBQP $\rightarrow$ "efficient" circuits for permanent

Best known algorithm for permanent (Ryser b3): $U\left(n 2^{n}\right)$

## Secure Classical Access to Quantum Cloud



## Secure Quantum access to Quantum Cloud

## Secure Quantum access to Quantum Cloud

Limited QClient

## Secure Quantum access to Quantum Cloud



Limited QClient

## Secure Quantum access to Quantum Cloud



Quantum Links
Limited QClient

## Secure Quantum access to Quantum Cloud



Limited QClient
Quantum Links Untrusted Server

## Secure Quantum access to Quantum Cloud



## Secure Quantum access to Quantum Cloud



Broadbent, Fitzsimons, and Kashefi 2009 - Universal Blind Quantum Computing Informational security

## Secure Quantum access to Quantum Cloud

QKD for encoding


Broadbent, Fitzsimons, and Kashefi 2009 - Universal Blind Quantum Computing Informational security

## Secure Quantum access to Quantum Cloud



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## Secure Quantum access to Quantum Cloud

QKD for encoding
Teleportation for computing


Limited QClient


Testing for verification


Broadbent, Fitzsimons, and Kashefi 2009 - Universal Blind Quantum Computing Informational security

## Computing with Teleportation

$$
J(\alpha):=\frac{1}{\sqrt{2}}\left(\begin{array}{cc}
1 & e^{i \alpha} \\
1 & -e^{i \alpha}
\end{array}\right)
$$

## Computing with Teleportation

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gate teleportation


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$$

gate teleportation

$$
\left.\left(\begin{array}{cccc}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & -1
\end{array}\right)+\ldots\right\rangle
$$

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Quantum Computer

Single qubit rotation


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Single qubit rotation
Quantum Computer

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Quantum Computer

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$$



Hiding the measurement result

## Gates Composition



## Gates Composition



## Gates Composition



Client-Server interactions

Re-writing


Re-writing


## Universal Blind Quantum Computings

$$
X=\left(\tilde{U},\left\{\phi_{x, y}\right\}\right)
$$

## Universal Blind Quantum Computings

$$
\begin{aligned}
& X=\left(\tilde{U},\left\{\phi_{x, y}\right\}\right)
\end{aligned}
$$

random single qubit generator

$$
\begin{gathered}
1 / \sqrt{2}\left(|0\rangle+e^{i \theta}|1\rangle\right) \\
\theta=0, \pi / 4,2 \pi / 4, \ldots, 7 \pi / 4
\end{gathered}
$$

## Universal Blind Quantum Computings


random single qubit generator

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## Universal Blind Quantum Computings


$\theta=0, \pi / 4,2 \pi / 4, \ldots, 7 \pi / 4$


## Universal Blind Quantum Computings



$$
\theta=0, \pi / 4,2 \pi / 4, \ldots, 7 \pi / 4
$$

$$
\begin{gathered}
r_{x, y} \in_{R}\{0,1\} \\
\delta_{x, y}=\phi_{x, y}^{\prime}+\theta_{x, y}+\pi r_{x, y}
\end{gathered}
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## Universal Blind Quantum Computings


random single qubit generator

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\theta=0, \pi / 4,2 \pi / 4, \ldots, 7 \pi / 4 \\
\underbrace{\delta_{x, y}=\phi_{x, y}^{\prime}+\theta_{x, y}+\pi r_{x, y}}_{r_{x, y} \in_{R}\{0,1\}}
\end{gathered}
$$



## Universal Blind Quantum Computings


random single qubit generator

$\left\{\left\{\left|+\delta_{x, x y}\right\rangle,\left|-\delta_{x_{x, y},}\right\rangle\right\}\right.$

## Universal Blind Quantum Computings


random single qubit generator

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\begin{gathered}
1 / \sqrt{2}\left(|0\rangle+e^{i \theta}|1\rangle\right) \\
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r_{x, y} \in_{R}\{0,1\} \\
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\end{gathered}
$$



$$
s_{x, y} \in\{0,1\} \quad\left\{\left|+\delta_{x, y}\right\rangle,\left|-\delta_{x, y}\right\rangle\right\}
$$

## Universal Blind Quantum Computings


random single qubit generator

$$
\begin{gathered}
1 / \sqrt{2}\left(|0\rangle+e^{i \theta}|1\rangle\right) \\
\theta=0, \pi / 4,2 \pi / 4, \ldots, 7 \pi / 4 \\
\underbrace{\in_{R}\{0,1\}}_{x, y} \\
\delta_{x, y}=\phi_{x, y}^{\prime}+\theta_{x, y}+\pi r_{x, y}
\end{gathered}
$$



$$
s_{x, y}:=s_{x, y}+r_{x, y} \quad s_{x, y} \in\{0,1\} \quad\left\{\left|+\delta_{x, y}\right\rangle,\left|-\delta_{x, y}\right\rangle\right\}
$$

## Security Definition

$$
\text { Protocol } \mathrm{P} \text { on input } X=\left(\tilde{U},\left\{\phi_{x, y}\right\}\right) \text { leaks at most } L(X)
$$

$\Rightarrow$ The distribution of the classical information obtained by Server is independent of $X$
$\Rightarrow$ Given the above distribution, the quantum state is fixed and independent of $X$

## What about correctness?

## What about correctness ?

- Correctness: in the absence of any deviation, client accepts and the output is correct
- Soundness: Client rejects an incorrect output, except with probability at most exponentially small in the security parameter


## Verification of Quantum Computing

## Self Testing 2005

Decide if the physical devices simulate their specification


## Verification of Quantum Computing

## Single-prover prepare-and-send

verifier has the ability to prepare quantum states and send them to the prover

- State authentication-based protocols
- Trapification-based protocols
- Test or Compute

| Protocol | Verifier resources | Communication | 2-way quantum comm. |
| :--- | :---: | :---: | :---: |
| Clifford-QAS VQC | $O(\log (1 / \epsilon))$ | $O(N \cdot \log (1 / \epsilon))$ | Y |
| Poly-QAS VQC | $O(\log (1 / \epsilon))$ | $O((n+L) \cdot \log (1 / \epsilon))$ | N |
| VUBQC | $O(1)$ | $O(N \cdot \log (1 / \epsilon))$ | N |
| Test-or-Compute | $O(1)$ | $O((n+T) \cdot \log (1 / \epsilon))$ | N |

## Verification of Quantum Computing

## Single-prover receive-and-measure

verifier receives quantum states from the prover and has the ability to measure them

- Post-hoc Verification (none hiding)
- Measuring only blind QC

| Protocol | Measurements | Observables | Blind |
| :--- | :---: | :---: | :---: |
| Measurement-only | $O\left(N \cdot 1 / \alpha \cdot 1 / \epsilon^{2}\right)$ | 5 | Y |
| Hypergraph measurement-only | $O\left(\max \left(N, 1 / \epsilon^{2}\right)^{22}\right)$ | 3 | Y |
| 1S-Post-hoc | $O\left(N^{2} \cdot \log (1 / \epsilon)\right)$ | 2 | N |
| Steering-based VUBQC | $O\left(N^{13} \log (N) \cdot \log (1 / \epsilon)\right)$ | 5 | Y |

## Verification of Quantum Computing

## Multi-prover entanglement-based

Classical Verifier interacts with more than one provers that are not allowed to communicate during the protocol

- CHSH game Rigidity
- Self-testing graph states
- Pauli Braiding

| Protocol | Provers | Qmem provers | Rounds | Communication | Blind |
| :--- | :---: | :---: | :---: | :---: | :---: |
| RUV | 2 | 2 | $O\left(N^{8192} \cdot \log (1 / \epsilon)\right)$ | $O\left(N^{8192} \cdot \log (1 / \epsilon)\right)$ | Y |
| McKague | $O\left(N^{22} \cdot \log (1 / \epsilon)\right)$ | 0 | $O\left(N^{22} \cdot \log (1 / \epsilon)\right)$ | $O\left(N^{22} \cdot \log (1 / \epsilon)\right)$ | Y |
| GKW | 2 | 1 | $O\left(N^{2048} \cdot \log (1 / \epsilon)\right)$ | $O\left(N^{2048} \cdot \log (1 / \epsilon)\right)$ | Y |
| HPDF | $O\left(N^{4} \log (N) \cdot \log (1 / \epsilon)\right)$ | $O(\log (1 / \epsilon))$ | $O\left(N^{4} \log (N) \cdot \log (1 / \epsilon)\right)$ | $O\left(N^{4} \log (N) \cdot \log (1 / \epsilon)\right)$ | Y |
| FH | 5 | 5 | $O\left(N^{16} \cdot \log (1 / \epsilon)\right)$ | $O\left(N^{19} \cdot \log (1 / \epsilon)\right)$ | N |
| NV | 7 | 7 | $O(1)$ | $O\left(N^{3} \cdot \log (1 / \epsilon)\right)$ | N |

## Verification of Quantum Computing

Overhead Noise
Scalability

## Trapification



Unconditionally Verifiable Blind Quantum Computing
Fitzsioms Kashefi, 2012

## Trapification



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Unconditionally Verifiable Blind Quantum Computing
Fitzsioms Kashefi, 2012

## Trapification

 Measurement angles remain hidden

Unconditionally Verifiable Blind Quantum Computing
Fitzsioms Kashefi, 2012

## Trapification

$\Omega_{\text {Eve,system }}$


## Trapification

## $\Omega_{\text {Eve,system }}$


$\sigma_{\text {testsubspace }}$

## Trapification



## Trapification



Prob trap being correct and the computation is wrong is bounded
$\sigma_{\text {testsubspace }}$

## Trapification



Prob trap being correct and the computation is wrong is bounded

$$
\begin{gathered}
\text { Trap Measurements } \\
M^{\theta}\left|+{ }_{\theta}\right\rangle \quad \rightarrow \quad s=0 \\
M^{\theta}\left|-{ }_{\theta}\right\rangle \quad \rightarrow \quad s=1
\end{gathered}
$$

$$
\sum_{\nu} p(\nu) \operatorname{Tr}\left(P_{\text {incorrect }}^{\nu} B(\nu)\right) \leq \epsilon
$$

$$
P_{\text {incorrect }}^{\nu}:=P_{\perp} \otimes|a c c\rangle\langle a c c|
$$

## Robust Verifiable Secure Quantum Access to Noisy Quantum Qloud



Classical input/output
Perfect blindness and exponential verification
Exponential correctness on honest-but-noisy device No overhead besides repetitions


Securing Quantum Computations in the NISQ Era

## Robust Verifiable Secure Quantum Access to Noisy Quantum Qloud



## Practical Efficient Honest Client - Malicious Server

Classical input/output
Perfect blindness and exponential verification
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Securing Quantum Computations in the NISQ Era

Kashefi, Leichtle, Music, Ollivier, 2020

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Computationally Secure (Post-quantum safe) Classical Access to Quantum Cloud?

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Classical Client Quantum FHE Mahadev, FOCS 2018

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Delegated Pseudo-Secret Random Qubit Generator
Cojocaru, Colisson, Kashefi, Wallden, AsiaCrypt 2019

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$O(1000)$ server qubits for randomising one single client qubits

$$
\begin{gathered}
r_{x, y} \in_{R}\{0,1\} \\
\delta_{x, y}=\phi_{x, y}^{\prime}+\theta_{x, y}+\pi r_{x, y}
\end{gathered}
$$





Secure Access to Quantum Cloud $\equiv$

Quantum Communication

## Malicious Client - Malicious Server



## Yao Garbled Circuit - Secure 2-party Computing



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## Yao Garbled Circuit - Secure 2-party Computing



Computational Security
Requires OT
Insert secret input b Evaluate f(a,b)
Honest but Curious Adversary

## Verifiable Quantum Yao



## Secret input q_c

Garbled CP map


## Verifiable Quantum Yao



## Verifiable Quantum Yao



## Verifiable Quantum Yao



## Secret input q_c

Garbled CP map


## Verifiable Quantum Yao



Server's input placed in DT(G) with correspoding trap-colouring


## Verifiable Quantum Yao



## Verifiable Quantum Yao



## Verifiable Quantum Yao



## Verifiable Quantum Yao



# Boosting Security <br> (Semi-Malicious Client to Fully Malicious one) 

Cut : Sender sends multiple copies of a state and message (with independent randomness) to the Receiver

## Practical Efficient Malicious Client - Malicious Server

states) wniere comectry comstucteu dy askmy tre oemuer io semu proofs and measuring them accordingly

Conditions for applying Q-CC $\longrightarrow$ Client manipulates single qubit


## Malicious Clients - Malicious Server



## Multiparty Delegated Quantum Computing 2017

Secret input q_1

## Garbled her part of the CP map



Secret input q_n
Garbled her part of the CP map


## Multiparty Delegated Quantum Computing 2017

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## Multiparty Delegated Quantum Computing 2017

Secret input q_1

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.


Secret input q_n
Garbled her part of the CP map

$$
\theta_{j}=\theta_{j}^{j}+\sum_{k=1, k \neq j}^{n}(-1)^{\oplus_{i=k}^{n} t_{j}^{j}} \theta_{j}^{k}
$$

## Multiparty Delegated Quantum Computing 2017

Secret input q_1

## Garbled her part of the CP map



Secret input q_n
Garbled her part of the CP map


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Secret input q_n
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Garbled her part of the CP map


## Multiparty Delegated Quantum Computing 2017

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Clients can insert traps only in their subgraphs

## But

A connected path for computation can be obtained only if they collaborate

## But

They need not to leak the position of traps

## Multiparty Delegated Quantum Computing 2017

Clients can insert traps only in their subgraphs

## But

A connected path for computation can be obtained only if they collaborate

## But

They need not to leak the position of traps


In Symmetric Case these issues are resolved by

## Multiparty Delegated Quantum Computing 2021

Double Blind QC - a classically orchestrated delegation

Good Enough State - correct up to a deviation independent of the inputs and security parameters


## VUBQC Deconstruction - Reconstruction

Steps to be updated to transform into a multi-client setting \&
Conditions that these replacement need to satisfy

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## Replacing Classical Steps with Classical SMPC



## Replacing Classical Steps with Classical SMPC



## Double Blind QC



## Double Blind QC



## Double Blind QC

Realised itself by a UBQC pattern


## Double Blind Gadgets for H or I



## Double Blind Gadgets for H or I



Clients: sends encrypted input and rotated states

SMPC: redistribute them to become dummy or trap

## Multiparty Delegated Quantum Computing 2021

Dulek, Grilo, Jeffery, Majenz, Schaffner 2020
Alon, Chung, Chung, Huang, Lee, Shen

| Metric | 9] | 26] | 1] | This work |
| :---: | :---: | :---: | :---: | :---: |
| Type | Stat. upgrade of CSMPC | Statistical | Comp. (FHE + CSMPC) | Stat. upgrade of CSMPC |
| Abort | Unanimous | Unanimous | Identifiable | Unanimous |
| Composability | Composable | Stand-Alone | Stand-Alone | Composable |
| Max Malicious Players | $N-1$ | $\left.\frac{C_{\text {dist }}-1}{2}\right\rfloor$ | $N-1$ | $N-1$ |
| Protocol Nature | Symmetric | Symmetric | Semi-Delegated | Delegated |
| Network Topology | Q and C: Complete | Q and C: Complete | Q and C: Complete | Q: Star / C: Complete |
| Q Operations | F.T. Q. Comp | FT Q Comp | FT Q Comp | Cl.: Single Qubit Serv.: FT Q Comp |
| Classical SMPC | Clifford Computation, <br> Operations in $\mathbb{Z}_{2}$, CT | CT | Clifford Computation, FHE verification | Operations in $\mathbb{Z}_{8}, \mathbb{Z}_{2}, \mathrm{CT}$ |
| Rounds (C or CSMPC) | $\mathcal{O}(g+\eta(N+t))$ | $d+2$ | $\mathcal{O}(1)$ | $d+5$ |
| Rounds (Q) | $\begin{gathered} \text { Par.: } \mathcal{O}(N d) \\ \text { Seq.: } \mathcal{O}(N(N+t+c)) \end{gathered}$ | Par.: 3 (2 if C output) <br> Seq.: $\mathcal{O}\left(\eta^{2}(N+t)\right)$ | Par.: $\mathcal{O}\left(N^{4}\right)$ | Par.: 2 (1 if C output) <br> Seq.: $\mathcal{O}(\eta N d)$ |
| Size of Q Memory | Par.: $\left.\mathcal{O}\left(\eta^{2}(N+t)\right)\right)$ <br> Seq.: $\mathcal{O}\left(\eta^{2} N\right)$ | $\begin{aligned} & \text { Par.: } \mathcal{O}\left(\eta^{2} N(N+t)\right) \\ & \text { Seq.: } \mathcal{O}\left(N^{2}\right) \end{aligned}$ | Par.: $\mathcal{O}\left(t N^{9} \eta^{2}\right)$ | $\begin{aligned} & \text { Cl.: } 3 \text { (0 if C I\&O) } \\ & \text { Serv. (par.): } \mathcal{O}\left(\eta N^{2} d\right) \\ & \text { Serv. (seq.): } \mathcal{O}(\eta N d) \end{aligned}$ |

Lipinska, Ribeiro, Wehner 2020

## Practical Efficient Malicious Clients - Malicious Server?

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## Each Module Can be Optimised

- SMPC : angles evaluations and permutations
- Remote State Prep : Hardware Dependent
- Blind QC : Not every qubits being hidden
- Verifiable QC : No Need for dummies

Key component = Remote State Preparation

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## The Most Optimal Client-Server RSP



## Quantum Enclave - Remote State Rotation



## The Most Optimal Multi Party QSMPC

## Qline Architecture + Remote State Rotation + QSMPC



VeriQloud's fully connected quantum network with a single optical fibre


## A Secure New World



