MoDaS

Mob Data Sourcing

Asking the Right Questions in Crowd Data Sourcing

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Outline

- **Introduction to crowd (data) sourcing**
- Databases and crowds
- Declarative is good
- How to best use resources
- Conclusion

Ack: Some slides are borrowed (with permission) from the VLDB'11 tutorial [DFKK11].

Disclaimer: - Very high level
- More questions than answers
- Some nudity ☺

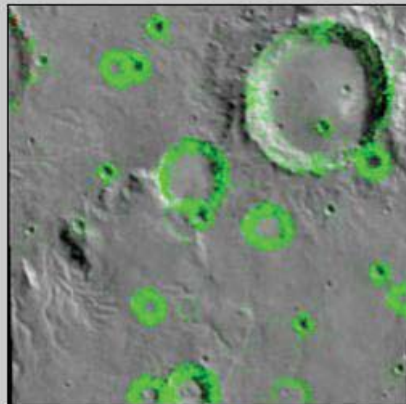# Crowd Sourcing 101

Billions of devices

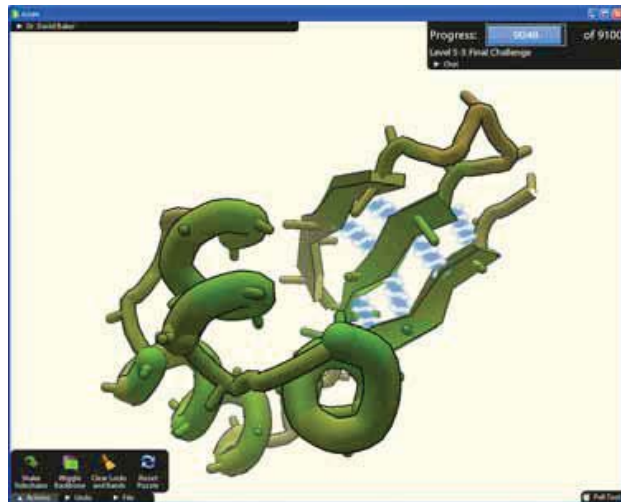# Crowd Sourcing 101

Ubiquitous connectivity

Examples

Citizen science

Pixels indicate Clickworker's identified craters

GALAXY ZOO — HUBBLE

Classify galaxies
Answer the question below using the buttons provided.

Is the galaxy simply smooth and rounded, with no sign of a disk?

Smooth   Features or disk   Star or artifact

foldit BETA
Solve Puzzles for Science

Examples

Citizen journalism and sensing

Examples

Examples

Games are fun!

So what is it all about?

- Bederson & Quinn (Human Computation) CHI'11
    - Motivation (Pay, altruism, enjoyment,…)
    - Quality control (we'll talk more about that)
    - Aggregation (We'll also talk more about that)
    - Human skills (Visual recognition, language, …)
    - …

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Databases and Crowds

- How can crowds help databases?
  - Fix broken data: entity resolution, inconsistencies
  - Add missing data
  - Subjective comparisons

- How can databases help crowd apps
  - Lazy data acquisition (only get the data that is needed)
  - Manage the data sourced from the crowd
  - Semi automatically create user interfaces

Database platforms for Crowd-based Data Sourcing

- Data models, query languages
  (query processing, optimization,…)
  - **Qurk (MIT)**
  - **CrowdDB (Berkley, ETH)**
  - **sCOOP (Stanford, UCSC)**
  - **FusionCOMP (TsuKuba)**
  - **MoDaS (Tel Aviv University)**

  - …

- Data quality

- Asking (the crowd) the right questions

Qurk (MIT)

- **Goal:** crowd-source comparisons, missing data
- **Basis:** SQL3 + UDF
  - UDF encapsulates crowd input
  - Special template language for crowd UDFs
  - Specify UI, quality control, possibly opt. hints

- **References:**
  [Marcus et al, CIDR'11, SIGMOD'11]

Qurk example

Is _____ Female?

men in a "people" database

ple(
(256),

TASK isFemale(tuple) TYPE:Filter
    Question: "is %s Female",
                Tuple["photo"]
    YesText: "Yes"
    NoText: "No"

e(p);

The magic is in the templates

- Templates generate UIs for different kinds of crowd-sourcing tasks
  - Filters: Yes/No questions
  - Joins: comparisons between two tuples (equality)
  - Order by: comparisons between two tuples (>=)
  - Generative: crowdsource attribute value

- Templates also specify quality control; e.g.
  COMBINER: MajorityVote

But, can you trust the Crowd?

Spencer Tunick

Many questions

- How to determine correctness ?

- How to clean the data?

- What questions to ask?

- Who to ask? (How many? When to stop?)

- How to best use resources?

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Example: Conflicts resolution

- Average value? Majority vote? Probabilistically?

- But some people know nothing about a given topic

- So maybe a "biased (probabilistic) vote"?

- But how to bias?

- A "chicken or the egg" problem:

    To know what is true we need to know who to believe.
    But to know this we need to know who is usually right
    (and in particular, what is true..)

Example: So what can we do?

- Start with some estimation on the trust in users

- Gain confidence in facts based on the opinion of users that supported them
  - Give bigger weight to user that we trust

- Then update the trust level in users, based on how many of the facts which they submitted, we believe

- Iterate until convergence

  Trusted users give us confidence in facts,
  and users that supported these facts gain our trust…

  [Galland et al, WSDM 2010]

- And there is also a probabilistic version…

But what do we want?

- Not yet another data cleaning algorithm

- We want to have easy control on the employed policy
  (for data cleaning, query selection, user game scores,...)

- We really don't want to (re)write Java code (for each tiny change!)

- We want (seamless) optimization, update propagation,...

Database approach:

　Define a **declarative language** for specifying policies

[Deutch, Greenshpan, Kostenko, M. ICDE'11 ,WWW'12]
[Deutch, Koch, M. PODS'10]

Proposed language

- Add to SQL (relational algebra) a REPAIR-KEY construct

    REPAIR-KEY "repairs" key violations in the database, choosing one possible option, probabilistically, according to the support

- And a WHILE construct

- Semantics: Markov chain of DB instances. Probability of a fact to hold in a given instance.

Name	Cuisine	support
Anton's	French	0.8
Anton's	Continental	0.2
McDonald	FastFood	1.0
...	...	

- Expresses nicely common policies for cleaning, selection of questions, scoring answers

TriviaMaster (ICDE 2011 demo)

Some complexity results

Formal problem: Given a Markov Chain of database instances and an SQL query on the database ("what is Anton's cuisine ?"), compute the probabilities of the different answers.

- <u>Theorem</u>: Exact computation is #P-hard

- <u>Theorem</u>: If Markov Chain is **ergodic**, computable in EXPTIME
  - Compute the stochastic matrix of transitions
  - Compute its fixpoint
  - For ergodic Markov Chain it corresponds to correct probabilities
  - Sum up probabilities of states where the query event holds

- <u>Theorem</u>: In general, 2-EXPTIME
  - Apply the above to each connected component of the Markov Chain
  - Factor by probability of being in each component

Some complexity (cont.)

Approximations:

– **Absolute approximation:** approximates correct probability ±ε
– **Relative approximation:** approximates correct probability up to a factor in-between (1- ε), (1+ ε).

[Relative is harder to achieve]

Language	Exact computation	Relative approx	Absolute approx
(Linear) datalog	#P-hard In PSPACE	NP-hard	In PTIME
Inflationary fixpoint	#P-hard In PSPACE	NP-hard	In PTIME
Non-inflationary fixpoint	#P-hard In (2)EXP-TIME	NP-hard	NP-hard; PTIME in input size and mixing time

Still lots of open questions

- How (and when) can we evaluate things fast enough?

- How to store the vast amount of data?
    - Distributed Databases? Map-reduce?

- The data keeps changing. How to handle updates?

- ...

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Partial knowledge

	q1	q2	q3	q4	q5	q6	...		
u1	a	5		b					
u2	a		3						
u3		5	3	b					
u4	b	2	3						
u5	c		3	a					
...									

- Goal: Compute an aggregate function f for each query, e.g.
  - Some metric of the distribution (e.g. entropy)
  - Most frequent answer
  - Aggregated value (e.g. average)

Increasing knowledge

- Limited overall resources

- Limited user availability

- Bounded resources per question

Which cells to resolve?

[Boim, Greenshpan, M., Novgorodov, Polyzotis, Tan. ICDE'12,...]

Quantifying uncertainty

- Assume t answers suffice for computing f for q

- Comp(q): all possible completions of q's column

- Dist(r – r'): distance between two results of f

- Uncertainty(q): max{ Dist(f(X) - f(Y)) | X,Y in Comp(q) }
  i.e. the largest distance between possibly completions

Quantifying uncertainty (cont.)

- Uncertainty measures for a Users-Answers matrix M
  - **Max-uncertainty(M)**
  - **Sum-uncertainty(M)**

- **Problem statement (X-uncertainty Reduction)**

  Given a matrix M, a choice $x \in \{max, sum\}$, and a set of constraints, identify a set C of empty cells that satisfy the constraints and where

  Max $_{M' \in M_C}$ X-uncertainty(M') is minimized.

  Where M_C contains all possible matrices that we can derive from M by resolving solely the cells in C.

Example

- ## Target function
  - Entropy, average, most frequent,…

- ## Constraints
  - A: bound k on the over number of cells
  - B: also a bound k' on questions per users
  - C: here k' is a bound on users per question

Some complexity results

- **max-Uncertainty Reduction**

    in PTIME for all constraints classes
    - Greedy algo for constraints class A (and C)
    - Using Max-flow for constraints class B

- **sum-Uncertainty Reduction**

    in PTIME for constraint classes A and C
    - Dynamic programming

    NP-COMPLETE for constraints class B
    - Reduction for perfect 3 set cover

AskIt (ICDE'12 demo)

- Gather information (scientific as well as fun)
  on ICDE'12 authors, participants, papers, presentations,…

Lots of open questions

- Use prior knowledge about users/answers
  - Predict answers
  - Predict who can/will answer what

  [Collaborative Filtering-style analysis is useful here]

- Worse-case analysis vs. expected error

- Treat other goal functions

- Optimization

- Incremental computation

…

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Conclusion

- All classical issues:

  Data models, query languages, query processing, optimization, HCI

- BUT
  - (Very) interactive computation
  - (Very) large scale data
  - (Very) little control on quality/reliability
  - Closed vs. open world assumption

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Mob Data Sourcing

‏תודה!‏

Thanks!

Merci!