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COMMUNICATIONS ET CRYPTOGRAPHIE AVEC DES VARIABLES QUANTIQUES CONTINUES

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Plan de l'exposé





* Les variables quantiques continues. des mesures Quantique Non Destructives à la cryptographie quantique...

* Cryptographie quantique avec des états cohérents (Nature 2003). du théorème de Shannon aux preuves de sécurité inconditionnelle

* Manipulation d'états non-gaussiens de la lumière (PRL 2004). de l'observation expérimentale d'états non-gaussiens à la distillation d'intrication et aux inégalités de Bell







* Essential feature : quantum channel with non-commuting quantum observables -> not restricted to single photon polarization !

-> New QKD protocol where :

* The non-commuting observables are the quadrature operators X and P

* The transmitted light contains weak coherent pulses (about 100 photons) with a gaussian modulation of amplitude and phase

* The detection is made using shot-noise limited homodyne detection







Initially introduced as a « criterion » for a « QND measurement » of X



Example 1 Fundamental idea for quantum key distribution: Alice and Bob encode information on X and P (and don't tell it in advance !)

Then $N_{eqB, X} = N_{eqB, P} = N_{eqB}$ and the best choice for Eve is $N_{eqE, X} = N_{eqE, P} = N_{eqE}$

Since everything is symmetric for X and P then :

 $N_{eqB} N_{eqE} \ge N_0^2$ (no-cloning theorem !)

(optimal cloning for QCV: $N_{eqB} = N_{eqE} = N_0$)





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Efficient transmission of information using continuous variables ? -> Shannon's formula (1948) : the mutual information I_{AB} (unit : bit / symbol) for a gaussian channel with additive noise is given by

 $I_{AB} = 1/2 \log_2 \left[1 + V(\text{signal}) / V(\text{noise}) \right]$

(a) Alice chooses X_A and P_A within two random gaussian distributions.

- (b) Alice sends to Bob the coherent state $|X_A + iP_A\rangle$
- (c) Bob measures either X_B or P_B
- (d) Bob and Alice agree on the basis choice (X or P), and keep the relevant values.





Main idea (Csiszar and Körner 1978, Maurer 1993) :

Alice and Bob can in principle distill, from their correlated key elements, a common secret key of size $S > sup(I_{AB} - I_{AE}, I_{AB} - I_{BE})$ bits per key element.

Crucial remark : it is enough that I_{AB} is larger than the **smallest** of I_{AE} and I_{BE} (i.e. one has to take the best possible case).



If I_{AE} is the smallest, the reconciliation must keep $S = I_{AB} - I_{AE}$ constant : Alice gives correction data to Bob (and also to Eve), and Bob orrects his data : « direct reconciliation protocol » If I_{BE} is the smallest, the reconciliation must keep $S = I_{AB} - I_{BE}$ constant : Bob gives correction data to Alice (and also to Eve), and Alice corrects his data : « reverse reconciliation protocol »





Reverse Reconciliation



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Bounding I_{BE} (F. Grosshans et al., *Nature* **421**, 238 (2003)) How well can Alice and Eve infer Bob's measurement results? Define the « conditional variance » $V(X_R | X_F) = V(X_R) - |\langle X_R X_F \rangle|^2 / V(X_F)$ Conditional variances are also bounded by Heisenberg relations : $V(X_B|X_A)_{min} V(P_B|P_E) \ge N_0^2$ $V(P_B|P_A)_{min} V(X_B|X_E) \ge N_0^2$ Using again Shannon's theorem... (and some algebra...) iff $T^{2}(N_{0} + N_{eaB})(N_{0} / V + N_{eaB}) < N_{0}^{2}$ $I_{BA} > (I_{BE})_{best}$

The security condition involves both T (channel transmission) and N_{eqB} (for direct reconciliation : $N_{eqB}\,<\,N_0$)



can be secure for any line transmission !







Coherent state QKD : experiment F. Grosshans et al., Nature **421**, 238 (2003)





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Example of exchanged data (burst of 60000 pulses @ 800 kHz, no on-line loss)









Secret key transmission : predicted results

Ideal SK rate : based on the error rate only, assumes perfect sofware efficiency (= ideal data extraction, reconciliation, and privacy amplification)

| | V _A | T _{line} | I _{BA} | I_{BE} (% of I_{BA}) | Ideal SK rate | Practical SK rate | |
|-------------------------|----------------|-------------------|-----------------|---------------------------|---------------------------------------|-------------------|--|
| | 40.7 | 1 | 2.39 | 0% | 1920 kb/s | | |
| | 37.6 | 0.79 | 2.17 | 58% | 730 kb/s | ? | |
| | 31.3 | 0.68 | 1.93 | 67% | 510 kb/s | | |
| | 26.0 | 0.49 | 1.66 | 72% | 370 kb/s | | |
| in shot- noise units | | | bits/ pulse | | Corresponding to a pulse rate 800 kHz | | |







At the end of the quantum exchange Alice and Bob share correlated strings of continuous data, from which they have to extract correlated bits.

Shannon's formula gives the maximum number of extractable bits, but this is an asymptotic value that requires adequate data processing ->

Optimized extraction method : "sliced reconciliation" :

N.J. Cerf, M. Lévy and G. Van Assche, PRA 63, 052311 (2001).







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Practical SK rate : final results, taking into account « all » imperfections

« Realistic » hypothesis : Eve cannot control the electronic noise floor of the homodyne detection (thermal noise in Bob 's resistors...)

| | V _A | T _{line} | I _{BA} | I_{BE} (% of I_{BA}) | Ideal SK rate | Practical SK rate | |
|---|-------------------------|-------------------|-----------------|---------------------------|---------------------------------------|-------------------|--|
| | 40.7 | 1 | 2.39 | 0% | 1920 kb/s | 1700 kb/s | |
| | 37.6 | 0.79 | 2.17 | 58% | 730 kb/s | 470 kb/s | |
| | 31.3 | 0.68 | 1.93 | 67% | 510 kb/s | 185 kb/s | |
| | 26.0 | 0.49 | 1.66 | 72% | 370 kb/s | 75 kb/s | |
| n | in shot- noise units | | bits/ pulse | | Corresponding to a pulse rate 800 kHz | | |







Cf BB84 vs entangled pair (Ekert) protocol

Unconditional security of coherent states QKD Institut d'Optique OIPC The DR and RR coherent states protocols are well within the Excess noise (N₀ units) separability « virtual entanglement » region ! entanglement Hint for unconditional security ? **Two approaches (both OK !):** RR + EPR states * Show that individual gaussian \mathbf{RR} + coherent states attacks are optimal -> secret rate ok! 3 9 6 (F. Grosshans, N. Cerf, PRL 2004) losses (dB)

DR : Direct Reconciliation RR : Reverse Reconciliation * Show that one can distill entangled qubits using CSS codes (S. Iblisdir, G. Van Assche, N. Cerf, PRL subm.)





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* Practical advantages of « actual » EPR beams vs. coherent states :

- The random values needed by Alice (encoding) and Bob (decoding) do not have to be externally generated (possibly by another quantum process), but they are produced by the protocol itself (« the key does not exist beforehand »).

- A « true » EPR protocol is more robust with respect to excess noise than a coherent state protocol (but the bit rates are the same if no excess noise).

- Marginal advantages ?

- * Fundamental advantage of « actual » EPR beams vs. coherent states :
- Entanglement distillation procedures and quantum repeaters !





















Conclusion





Security proof of coherent state QKD :

- * Coherent states protocols using reverse reconciliation are secure against any (gaussian or non-gaussian) finite-size attack
 * Unconditional security of these protocols has also been (almost) proven.
- Coherent states QKD demonstrator : Nature 421, 238 (2003)
- * Measured secure bit transmission rates : 1.7 Mbit/sec @ 0 dB loss 75 kbit/sec @ 3.1 dB loss
- * Competitive against faint pulses ? Test @ 1550 nm under way

Conditional preparation of « squeezed » non-gaussian pulses (PRL 2004)
* Phase-dependant non-gaussian Wigner function (« squeezed volcano »)
* First step towards : entanglement distillation procedures ?
new tests of Bell's inequalities ?